

Provable security of advanced properties of TLS and SSH

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Real World
Cryptography



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TLS vs. SSH

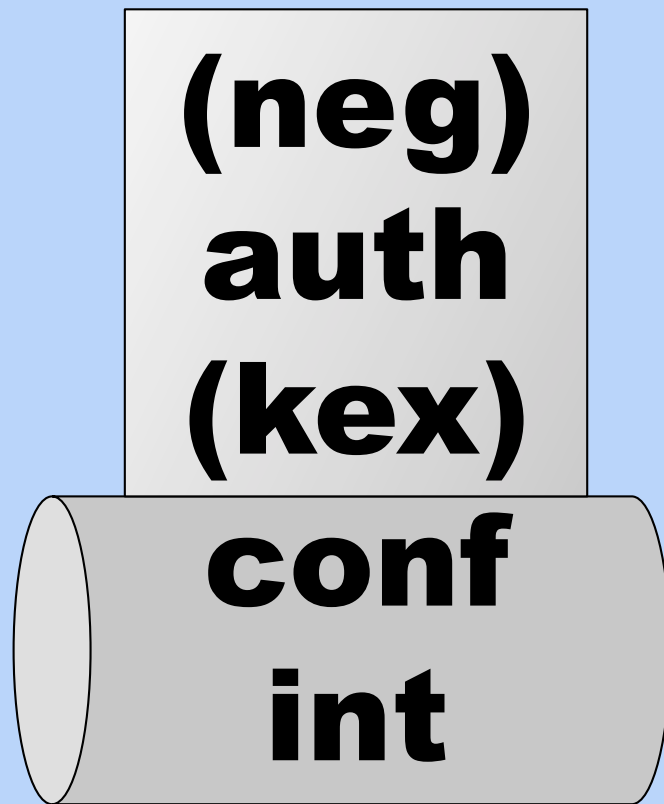
TLS

- provides secure transport for many applications
- entity authentication
- confidentiality & integrity of transmissions
- handshake establishes secure channel

SSH

- provides secure transport primarily for remote shell logins
- entity authentication
- confidentiality & integrity of transmissions
- handshake establishes secure channel

Security goals of TLS and SSH



From an application perspective, TLS and SSH provide:

- (negotiation of parameters)
- entity authentication
- (key exchange)
- confidentiality and integrity of messages

Security of TLS

Structure of TLS

HANDSHAKE PROTOCOL

Negotiation of cryptographic parameters

Authentication (one-way or mutual) using public key certificates

Establishment of a master secret key

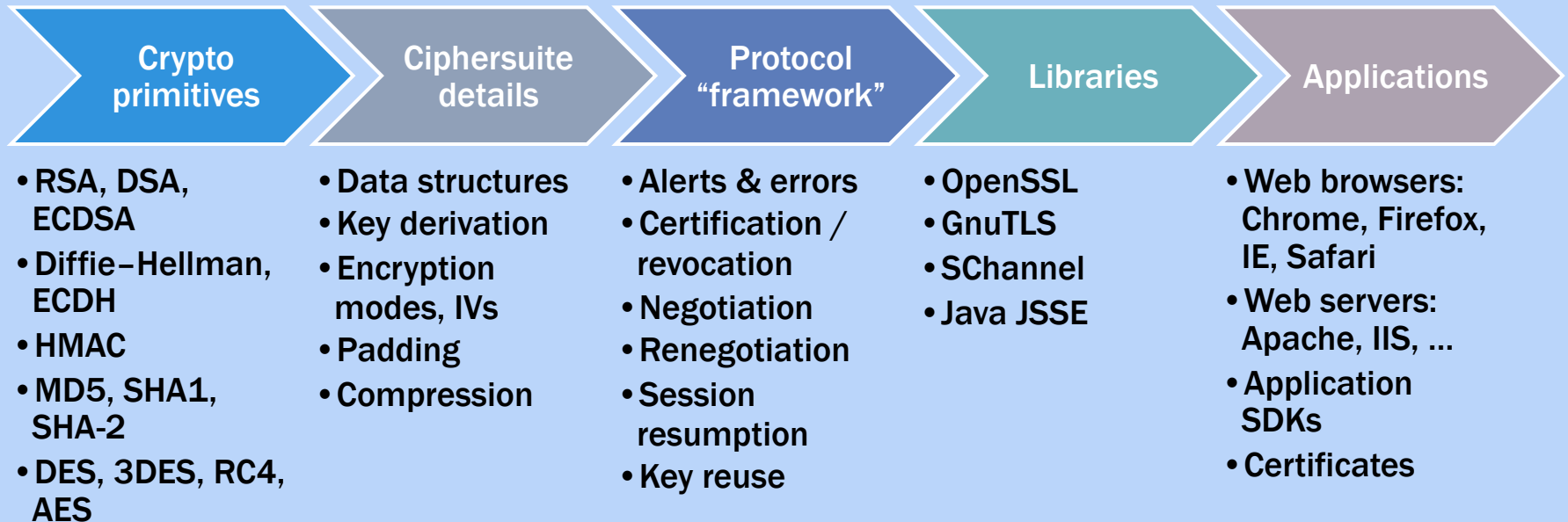
Derivation of encryption and authentication keys

Key confirmation

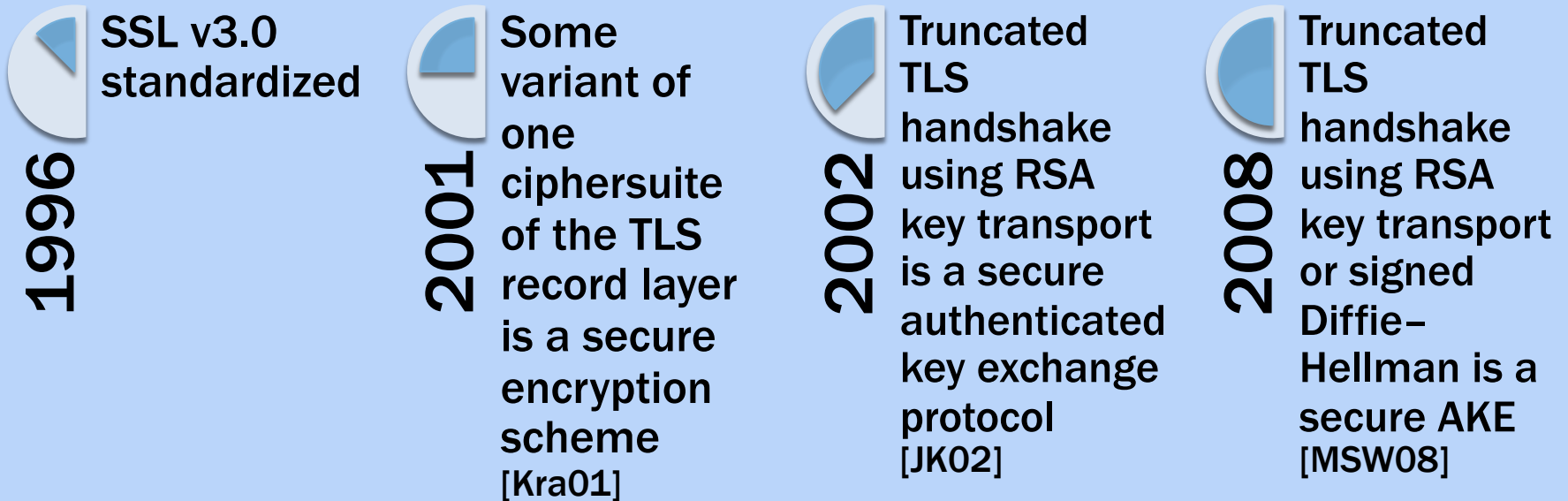
RECORD LAYER

Authenticated encryption of application data

Components of TLS

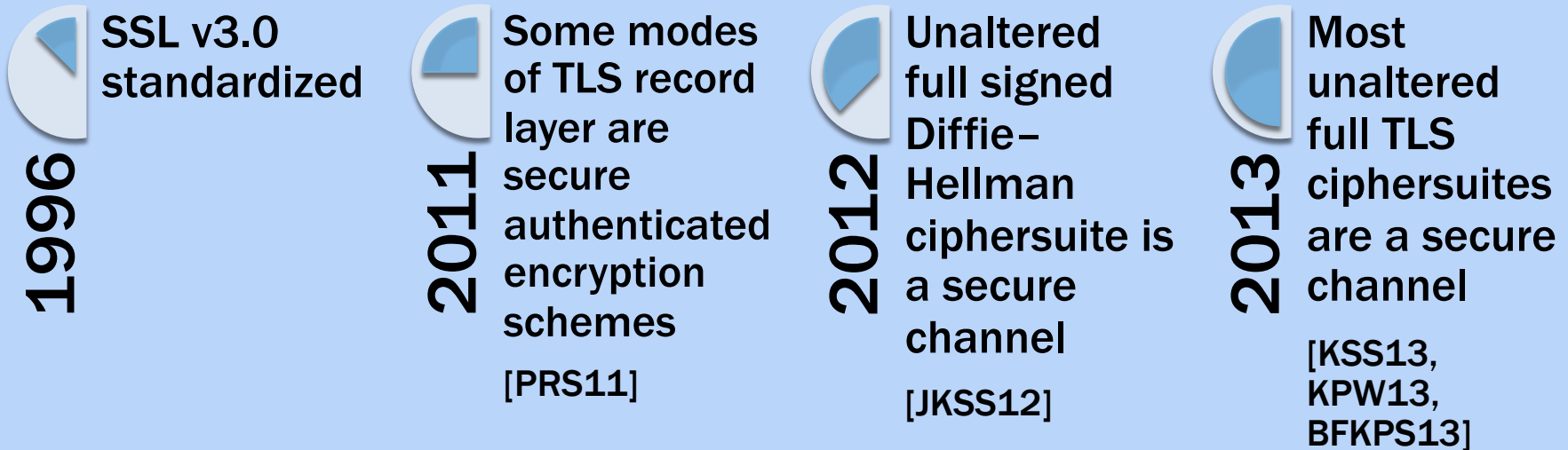


Is TLS secure?



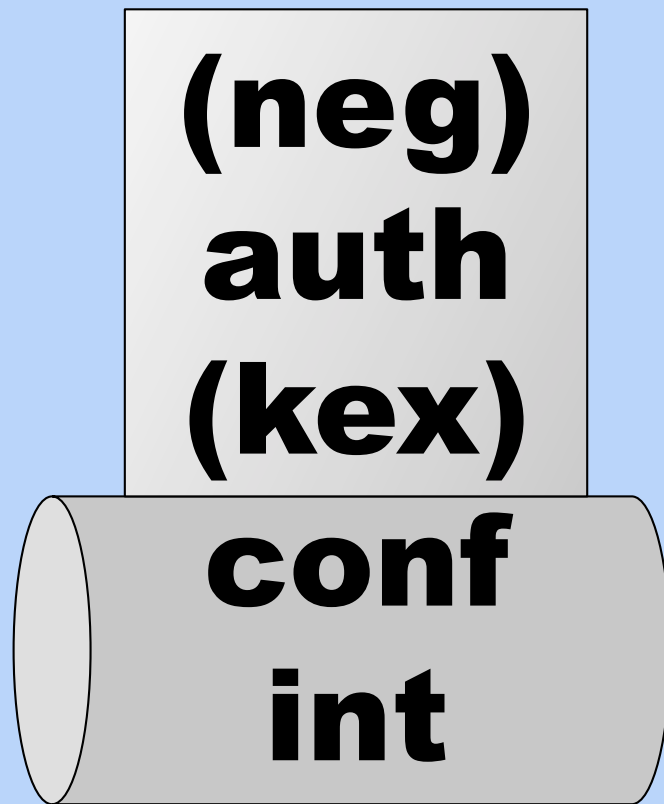
**“some variant” ... “truncated TLS” ...
limited ciphersuites**

Is TLS secure?



“unaltered” ... “full” ... “most ciphersuites”

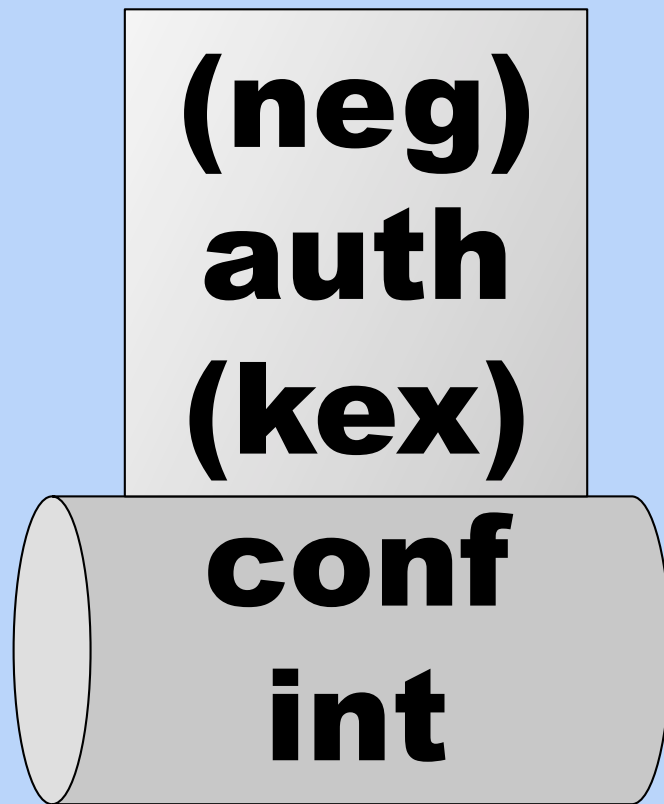
Security goals of TLS and SSH



From an application perspective, TLS and SSH provide:

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Security goals of TLS and SSH

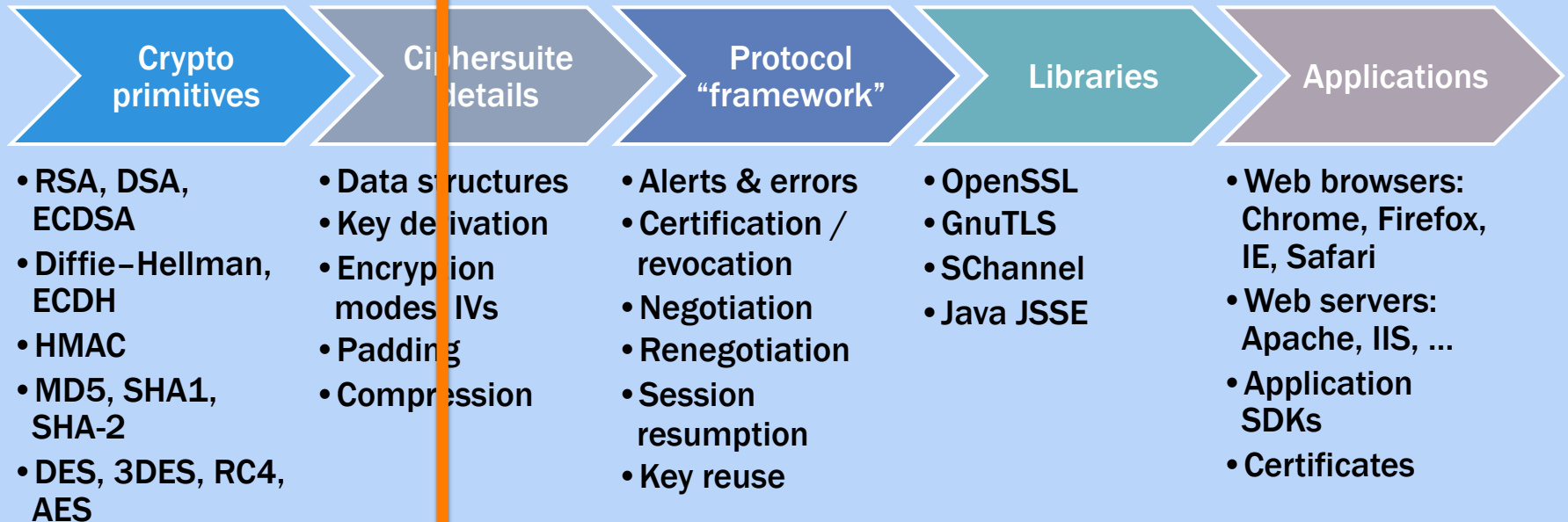


Authenticated and Confidential Channel Establishment (ACCE)

security definition
[JKSS12] captures:

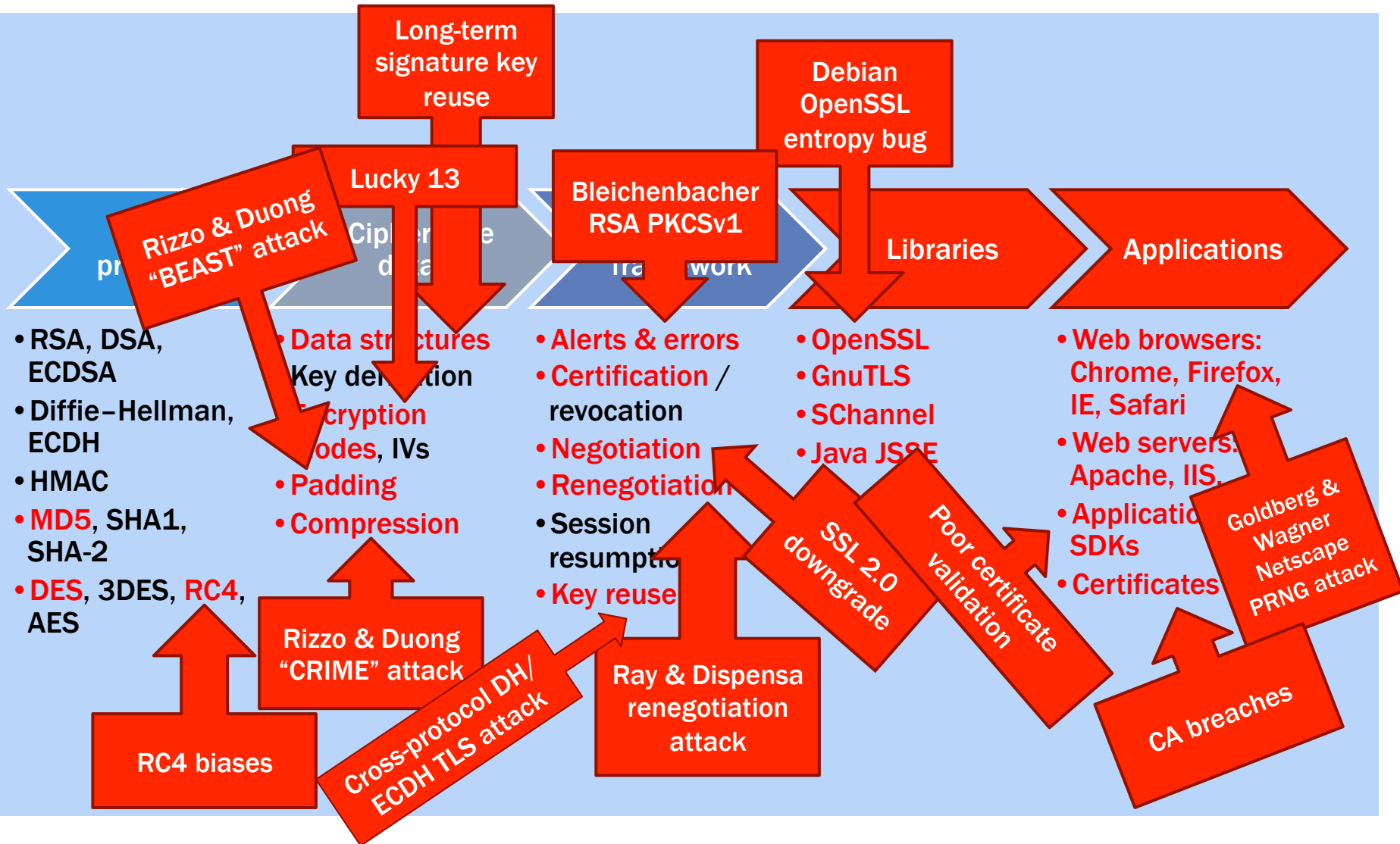
- entity authentication
- (key exchange)
- confidentiality and integrity of messages

Results on the provable security of TLS



Good combinations of these are theoretically secure, when done properly.

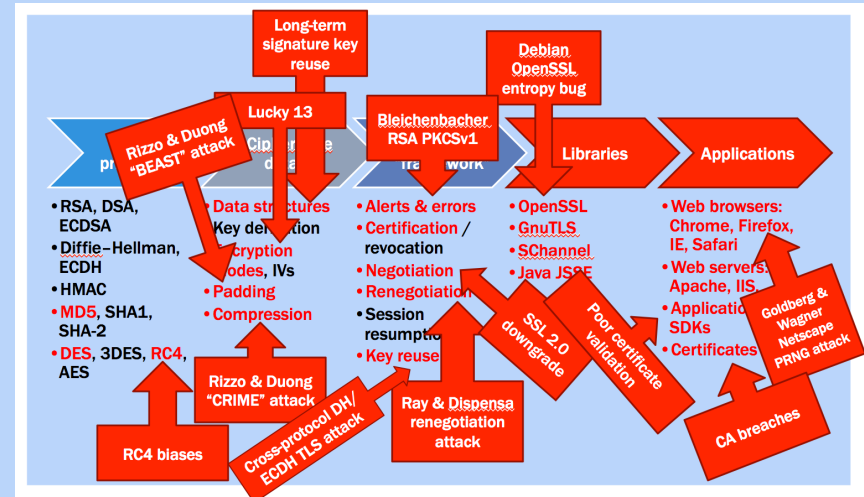
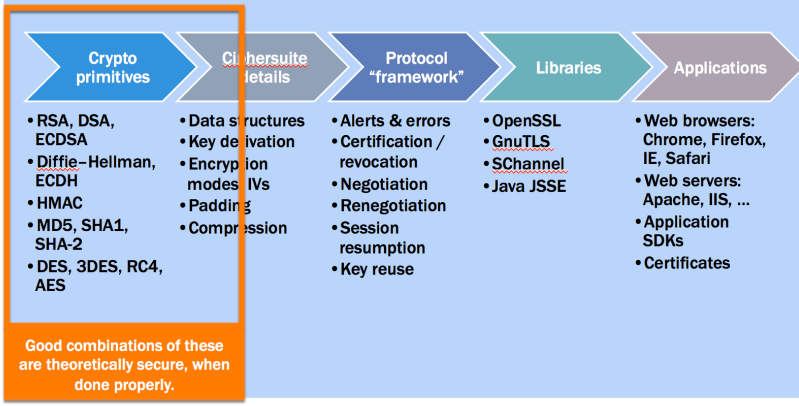
Real-world attacks on TLS



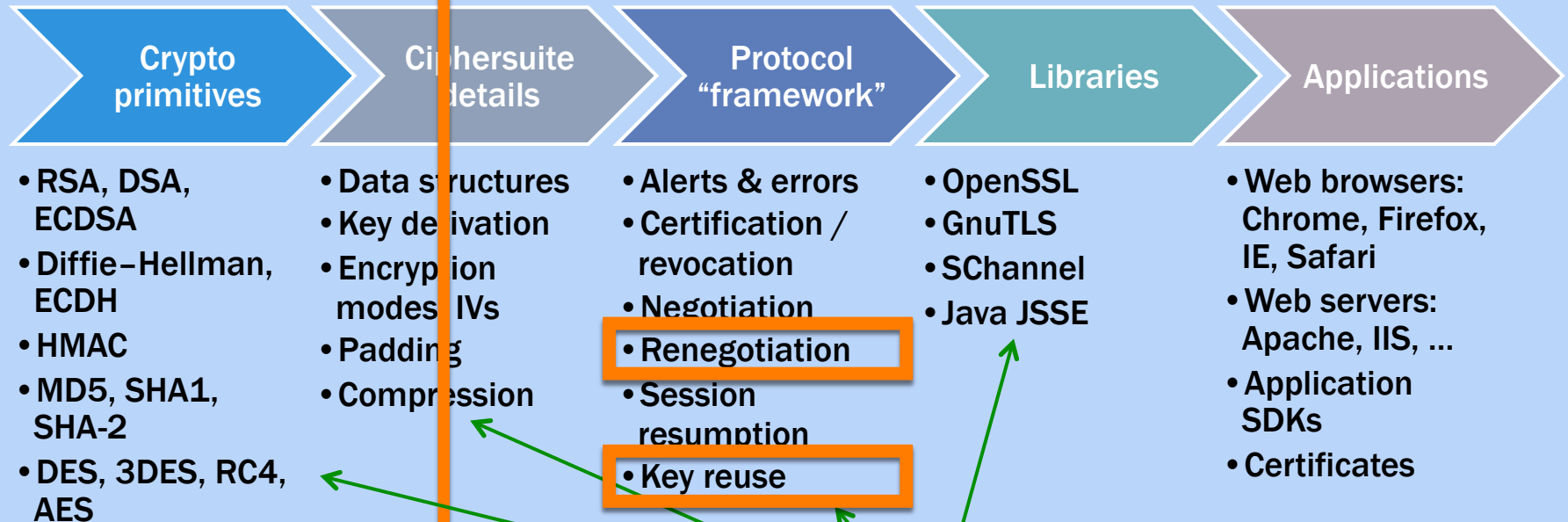
The gap between theory and practice

Provable security results

Real-world attacks



Components of TLS



Good combinations of these are theoretically secure, when done properly.

Other recent work [BFKPS13,BFKPSZ14] looks at several layers simultaneously.

TLS and renegotiation

ACM CCS
2013
eprint
2012/630

Why renegotiate?

Renegotiation allows parties in an established TLS channel to create a new TLS channel that continues from the existing one.

Once you've established a TLS channel, why would you ever want to renegotiate it?

- Change cryptographic parameters
- Refresh encryption keys (“more forward secrecy”)
- Change authentication credentials
- Identity hiding for client
 1. Establish a one-way authenticated TLS session
 2. Renegotiate using mutual authentication.
Since handshake messages are sent in the encrypted TLS channel, client's identity is kept private.

Renegotiation in TLS

(pre-November 2009)

Client

Server
(TLS)

TLS handshake₀

TLS recordlayer₀

m₀

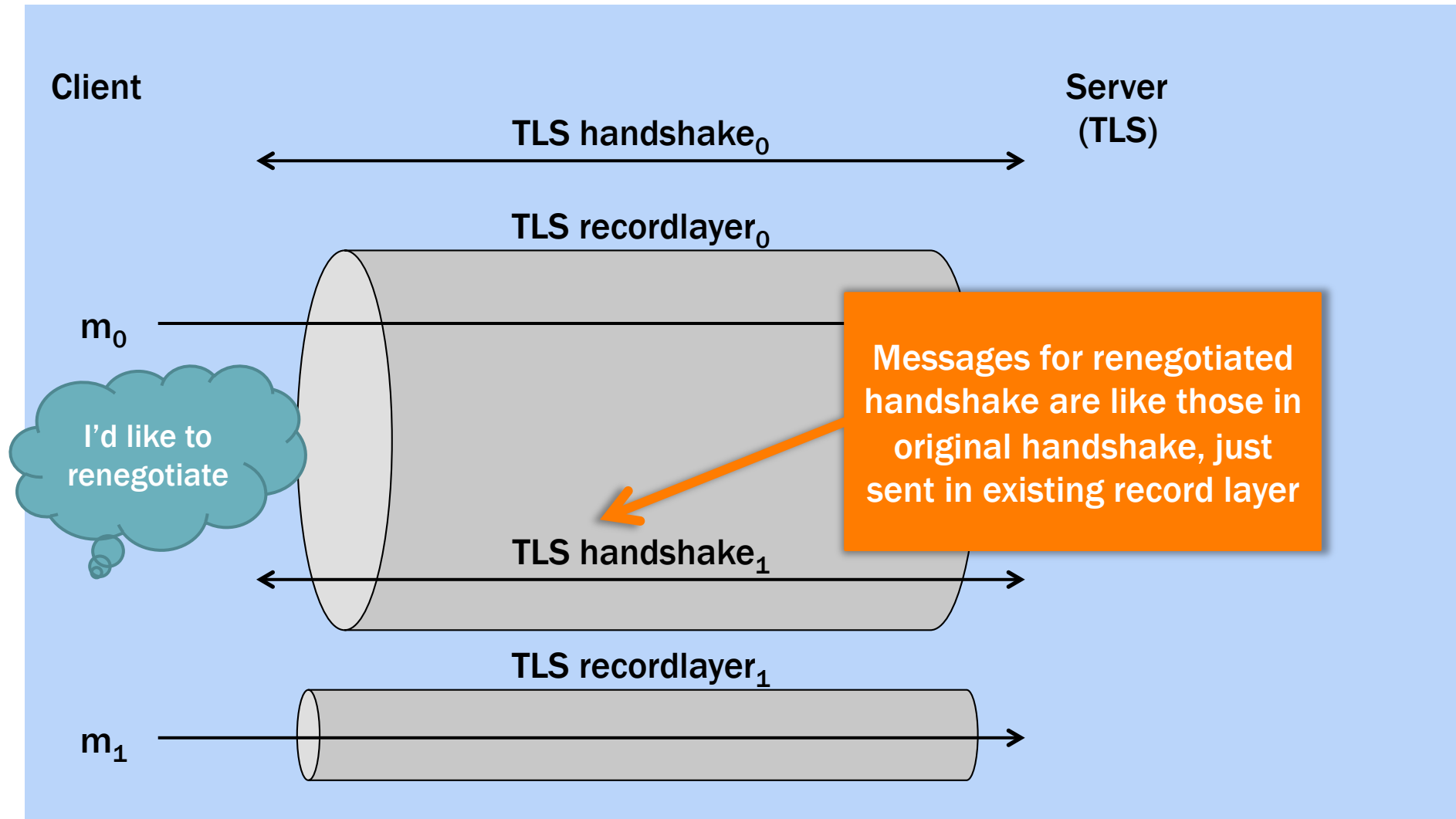
I'd like to renegotiate

Messages for renegotiated handshake are like those in original handshake, just sent in existing record layer

TLS handshake₁

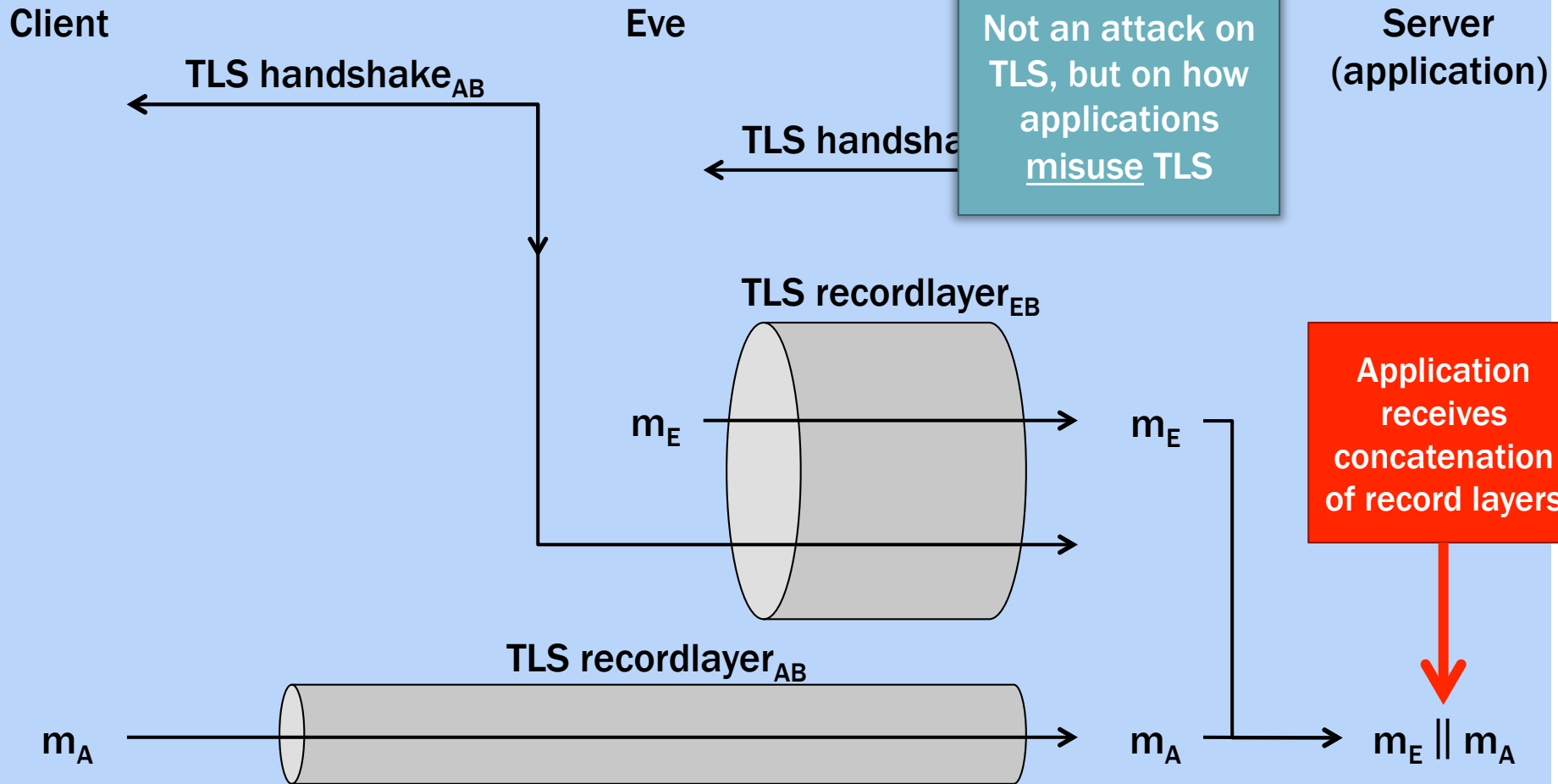
TLS recordlayer₁

m₁



TLS Renegotiation “Attack”

Ray & Dispensa, November 20



Renegotiation security

Q: What property should a secure renegotiable protocol have?

A: Whenever two parties successfully renegotiate, they are assured they have the exact same view of everything that happened previously.

TLS Renegotiation Countermeasures

Two related countermeasures standardized by IETF in RFC 5746:

1. Signalling Ciphersuite Value
2. Renegotiation Indication Extension

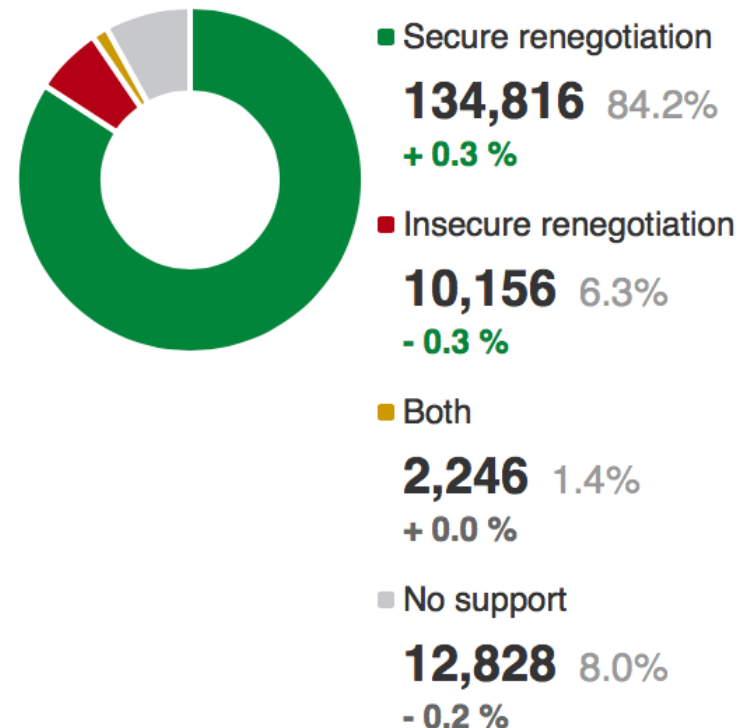
Basic idea: include fingerprint of previous handshake when renegotiating.

TLS Renegotiation Countermeasures

SCSV/RIE fairly quickly and widely adopted.

Currently 84% deployment
(SSL Pulse, Jan. 3, 2014)

Renegotiation Support



**Does this really fix the
problem?**

Does this really fix the problem?

ACCE security isn't enough: these ciphersuites have been proven ACCE security yet are vulnerable to renegotiation attack.

Need a security definition that includes renegotiation.

Technical approach

1. Extend authenticated and confidential channel establishment (ACCE) security model to include renegotiable, multi-phase protocols.
2. Define security notion for renegotiable protocols.
 - weakly secure renegotiable ACCE
 - secure renegotiable ACCE
3. Show that TLS without fixes does not satisfy security definition.
4. Show that TLS with fixes does satisfy security definition.
 - Generic reduction: If a TLS ciphersuite satisfies a certain property, then, when combined with fixes, it is a weakly secure renegotiable ACCE.
5. Propose even stronger fix.

Secure renegotiable ACCE

Definition

When a party successfully renegotiates a new phase, its partner has a phase with a matching handshake and record layer transcript

TLS

- TLS with or without RFC 5746 fixes is not a secure renegotiable ACCE.
- Can be fixed by including hash of previous record layer messages in RIE.

Weakly secure renegotiable ACCE

Definition

When a party successfully renegotiate a new phase, its partner has a phase with a matching handshake and record layer transcript, *provided no previous phase's session key was revealed.*

TLS

- TLS without fixes is not a weakly secure renegotiable ACCE.
- TLS with RFC 5746 fixes is a weakly secure renegotiable ACCE.
 - (This is good enough.)

Lessons learned: TLS renegotiation

Theory

- Renegotiation not previously included in AKE security definitions.
 - Different levels of renegotiation security
- Security of a protocol in isolation doesn't imply security with renegotiation.
- Need to “open up” security definitions in order to generically compose protocols.

Practice

- Confidence in standardized TLS renegotiation fixes.

Multi-ciphersuite security, TLS and SSH

eprint
2013/813

List of SSH ciphersuites

■ Authentication:

- RSA signatures
- DSA-SHA1
- ECDSA-SHA2
- X509-RSA signatures
- X509-DSA-SHA1
- X509-ECDSA-SHA2

■ Key exchange:

- DH explicit group SHA1
- DH explicit group SHA2
- DH group 1 SHA1
- DH group 14 SHA1
- ECDH-nistp256-SHA2
- ECDH-nistp384-SHA2
- ECDH-nistp521-SHA2
- ECDH-*-SHA2
- GSS-group1-SHA1-*
- GSS-group14-SHA1-*
- GSS explicit group SHA1
- RSA1024-SHA1
- RSA2048-SHA2
- ECMQV-*-SHA2

■ Encryption:

- 3des-cbc
- blowfish-cbc
- twofish256-cbc
- twofish-cbc
- twofish192-cbc
- twofish128-cbc
- aes256-cbc
- aes192-cbc
- aes128-cbc

- serpent256-cbc
- serpent192-cbc
- serpent128-cbc
- arcfour
- idea-cbc
- cast128-cbc
- des-cbc
- arcfour128
- arcfour256
- aes128-ctr
- aes192-ctr
- aes256-ctr
- 3des-ctr
- blowfish-ctr
- twofish128-ctr
- twofish192-ctr
- twofish256-ctr
- serpent128-ctr
- serpent192-ctr
- serpent256-ctr
- idea-ctr
- cast128-ctr
- AEAD_AES_128_GCM
- AEAD_AES_256_GCM

■ MACs:

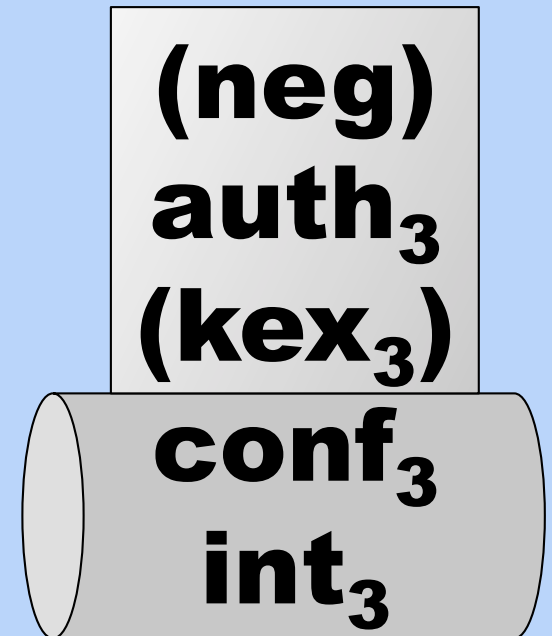
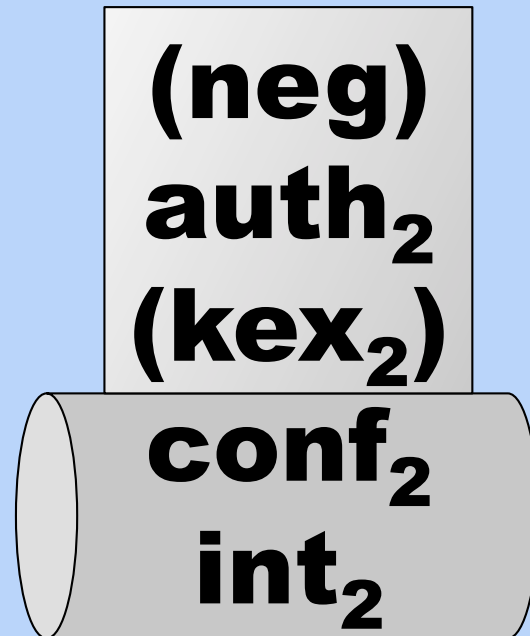
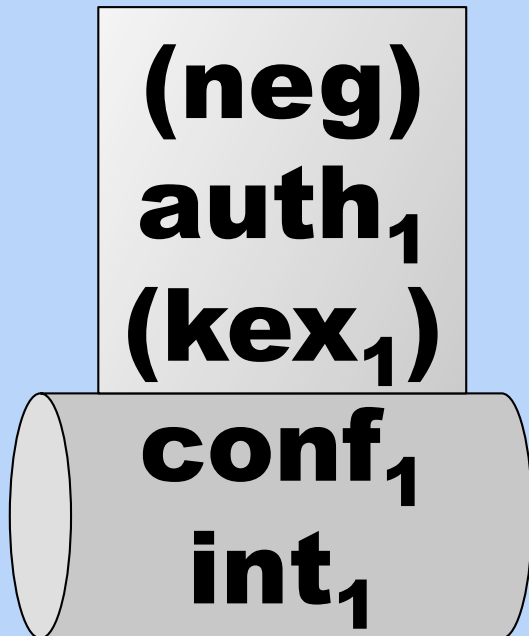
- hmac-sha1
- hmac-sha1-96
- hmac-md5
- hmac-md5-96
- AEAD_AES_128_GCM
- AEAD_AES_256_GCM
- hmac-sha2-256
- hmac-sha2-512

How we'd like to analyze ciphersuites

ciphersuite 1

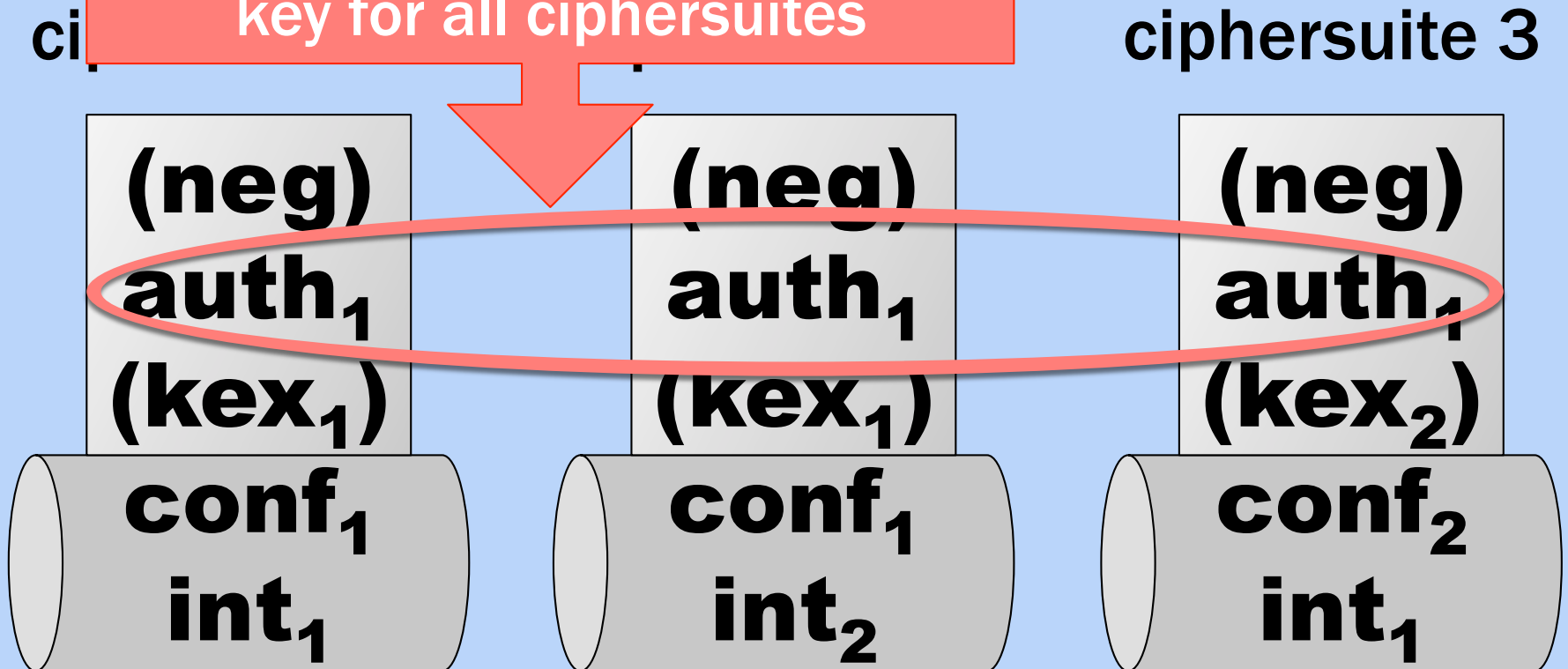
ciphersuite 2

ciphersuite 3



The reality of multi-ciphersuite usage

In practice, TLS and SSH servers use the same long-term key for all ciphersuites



Long-term key reuse across ciphersuites

Is this secure?

Even if a ciphersuite is provably secure on its own, it may not be secure if the long-term key is shared between two ciphersuites.

Long-term keys in TLS

Most TLS ciphersuites are provably secure channels (ACCE).

But this assumes that each ciphersuite uses its own distinct long-term key.

Long-term keys in TLS

In signed Diffie–Hellman ciphersuites, the thing that is signed is the raw DH parameters.

- i.e., without type information

Possible to interpret a set of ECDH params as valid finite field DH params.

- [MVVP12]

=> TLS not secure with long-term key reuse.

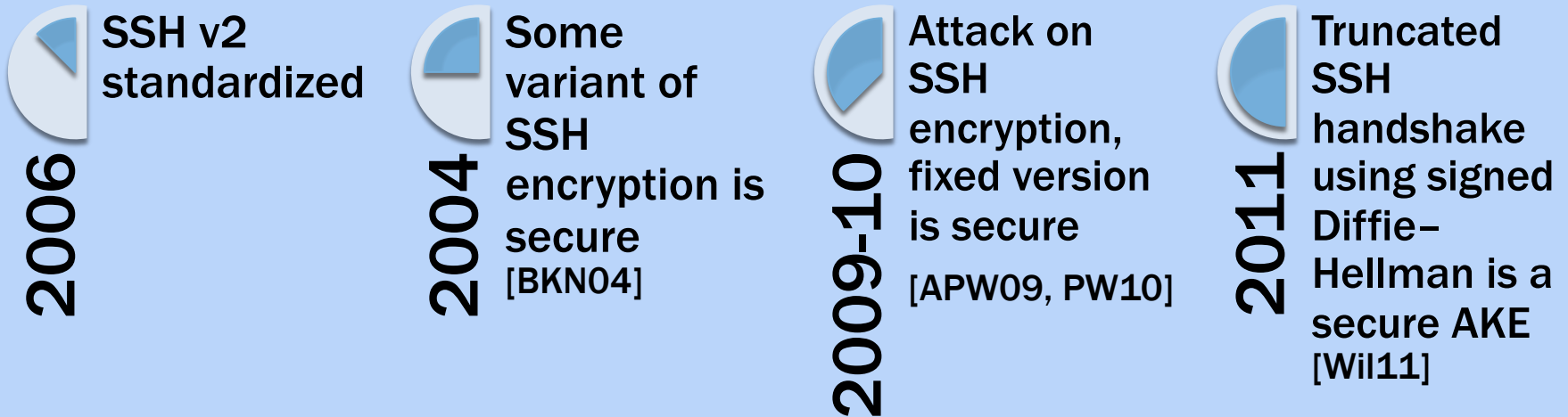
ACCE security of a ciphersuite in isolation does not imply security with long-term key reuse.

Long-term keys in SSH

In SSH, the thing that is signed contains an unambiguous identification of the intended ciphersuite.

We might hope to be able to prove SSH secure even with key reuse across ciphersuites.

Is SSH secure?



“some variant” ... “truncated SSH”

Signed DH SSH is a secure ACCE

Theorem: Assuming the signature scheme is secure, the computational Diffie-Hellman problem is hard, and the encryption scheme is a secure buffered stateful authenticated encryption scheme, then signed-DH SSH is a secure ACCE protocol.

How can we prove it secure even with long-term key reuse across ciphersuites?

Provable security of long-term key reuse

Goal: Generic theorem: If an individual ciphersuite is secure, then it is secure even if long-term keys are reused across ciphersuites.

- Impossible: TLS attack from previous slide.

Revised goal: Generic theorem: If an individual ciphersuite is secure under additional conditions, then it is secure even if long-term keys are reused across ciphersuites.

Lessons learned: multi-ciphersuite

Theory

- Definition for security of multi-ciphersuite protocols.
- Generic theorem on when it is safe to reuse long-term keys across individually secure ciphersuites.
 - Main idea: adding an auxiliary “signing oracle” to individual security to enable reduction, parameterize freshness condition.
 - Lots of other applications of this main idea...

Practice

- Confidence in signed-DH SSH ciphersuites, even if the same long-term keys are reused across ciphersuites.
 - ... and even when reused with unambiguously independent protocols.

Summary

Theory

- Provable security of single ciphersuites in isolation doesn't imply security in complex settings:
 - TLS renegotiation attack
 - multi-ciphersuite security
- Can augment ACCE security models for more complex functionality
- By opening up ACCE security models, can prove more generic composition theorems

Practice

- Confidence in TLS standardized renegotiation fixes.
- Confidence in SSH signed-DH ciphersuites in isolation or with long-term key reuse.

Cryptographers and standards

Tools for cryptographic modelling were not fully developed when key standards (TLS, SSH) were first proposed.

- Some flaws: padding, encryption modes, ...
- Some successes: general structure adequate

Encourage cryptographers and standards community to engage to support development & adoption of protocols with formal security properties where possible.